

## Motivation

- partial reconfiguration enables time-sharing of reconfigurable hardware resources
- hardware threads, as implemented by ReconOS, provide partitioning of an application into suitable modules for hardware multitasking
- non-preemptive multitasking techniques are unsuitable for many applications
  - long-running threads may make system unresponsive
  - asynchronous (i.e. blocking) operations must be registered with an event loop via callback functions
- preemptive multitasking faces substantial challenges when applied to partially reconfigurable devices
  - determining and accessing the relevant context of a hardware module is a complex task
  - readback or scan chain techniques involve significant overheads and are often device-dependent

## ReconOS Programming Model

### similar to existing APIs

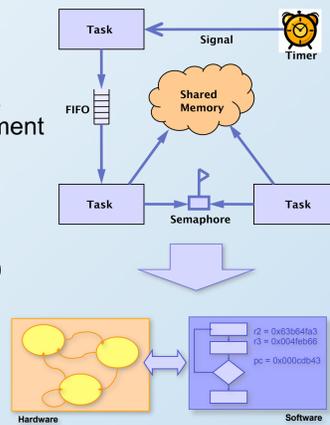
- eCos
- POSIX

### OS services

- task management
- memory management
- synchronization
- communication

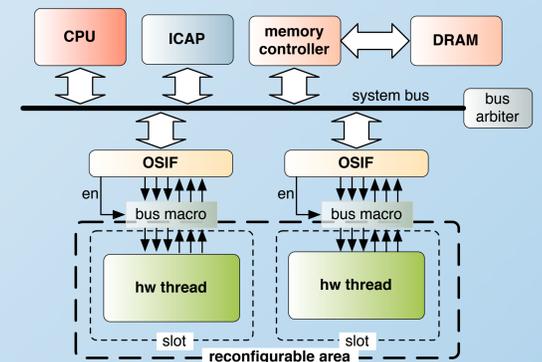
### OS objects

- tasks (HW or SW)
- shared memory
- semaphores
- queues/FIFOs
- timers
- signals
- ...



## ReconOS Execution Model

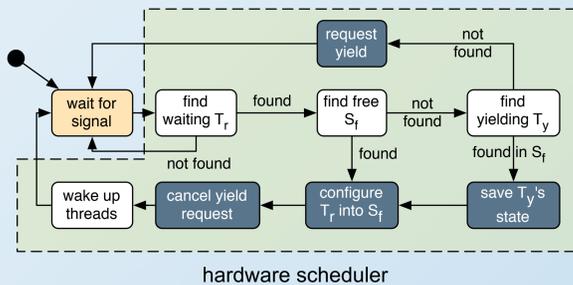
- OS interface module (OSIF) enables transparent communication and synchronization between hardware and software
- OS calls from hardware are relayed to *delegate threads* running on the system's CPU
- HW multitasking through partial reconfiguration



## Cooperative Multitasking

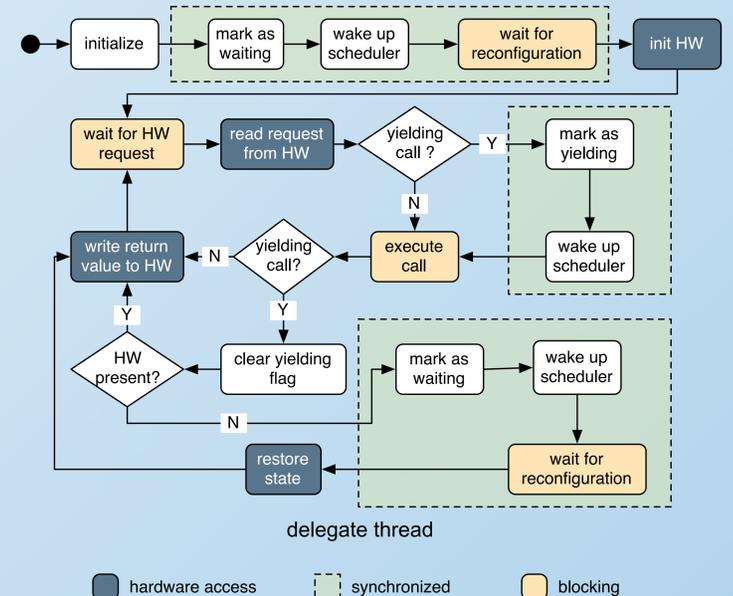
### Approach

- threads can voluntarily relinquish (`yield()`) their execution slot
- threads are responsible for saving and restoring their state on yield or resume
- ideally, threads yield on blocking OS calls, during which they would not perform any computations



### Implementation in ReconOS

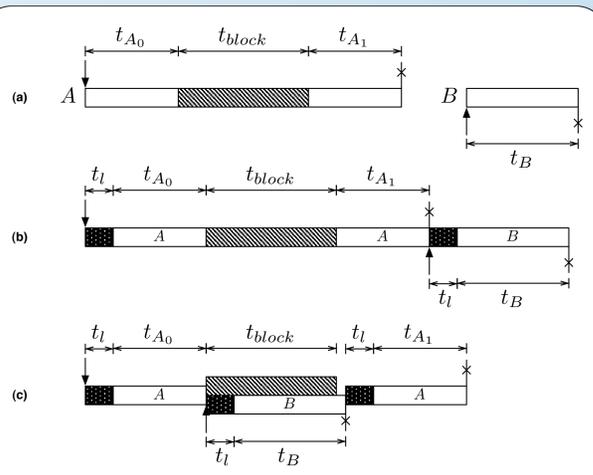
- in ReconOS, cooperative multitasking is only employed for HW threads; SW threads are scheduled preemptively
- the task of managing the reconfigurable resources is shared between two software threads
  - a hardware thread's *delegate* thread and a high-priority *hardware scheduler* thread
  - no changes to the OS kernel are necessary
- a synthesized hardware circuit representing a thread's functionality is called a *core*
- for every *slot* in the system, a core is placed and routed, resulting in  $n_{slots} \times n_{cores}$  partial *bitstreams*
- data structures modeling the relationships between slots, hardware threads, cores, and bitstreams are shared between the delegates and the hardware scheduler



hardware access (blue square), synchronized (green dashed box), blocking (yellow square)

## Scheduling Example

- (a) consider two threads, A and B
  - thread A runs for  $t_{A0}$ , blocks for  $t_{block}$ , and then runs again for  $t_{A1}$
  - thread B simply runs for  $t_B$
  - loading a thread onto the FPGA takes  $t_l$



- (b) with non-preemptive multitasking, threads A and B are executed consecutively, with a total run time

$$T_n(A, B) = 2t_l + t_{A0} + t_{A1} + t_{block} + t_B$$

- (c) with cooperative multitasking, thread A can *yield* its execution slot to thread B while blocking (i.e. on an OS call), resulting in an execution time of

$$T_c(A, B) = 2t_l + t_{A0} + t_{A1} + t_{As} + t_{Ar} + \max(t_{block}, t_l + t_B)$$

$t_{As}$  and  $t_{Ar}$  are the times to save and restore A's state, respectively.

- Thus, the cooperative multitasking approach reduces the total run-time, provided that both

$$t_B > t_{As} + t_{Ar} \quad \text{and} \quad t_{block} > t_{As} + t_{Ar} + t_l$$

## Outlook / Future Work

- efficient scheduling algorithms for a cooperatively multitasking subset of hardware threads in a preemptively scheduled multithreaded software system
- improved reconfiguration infrastructure to decrease reconfiguration overhead

## References

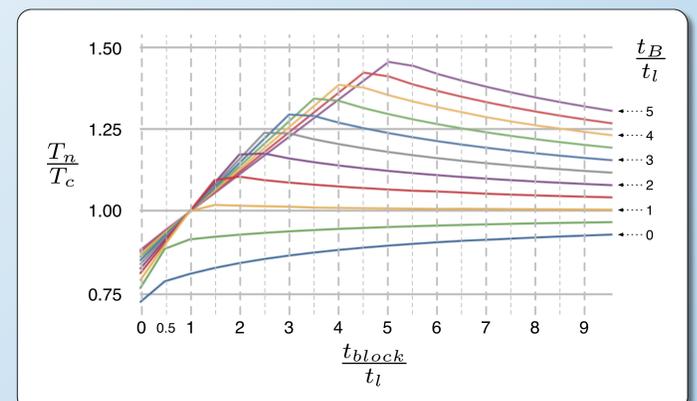
- E. Lübbers and M. Platzner, „Multithreaded Programming for Reconfigurable Computers,“ *ACM Transactions on Embedded Computing Systems (TECS)*, 2009, to appear
- E. Lübbers and M. Platzner, „ReconOS: An RTOS supporting Hard- and Software Threads,“ in *17th IEEE International Conference on Field Programmable Logic and Applications (FPL)*, 2007

## Experimental Results

- timing overheads of individual OS operations

thread initialization	1.76 ms
thread suspend	93.12 $\mu$ s
thread resume	192.32 $\mu$ s
state save (4096 bytes)	37.51 $\mu$ s (104.1 MB/s)
state restore (4096 bytes)	45.19 $\mu$ s (86.4 MB/s)
reconfiguration time (233 kBytes)	99.96 ms

- application execution time of a prototype implementation of the scheduling example



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